

#### Knapsack

0-1 Knapsack Example Time Complexit

Dynamic Programming

## Knapsack

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Dynamic Programming

## **0-1 Knapsack**

Given *n* objects with integer weights  $w_i$  and values  $v_i$ , what is the most valuable subset that weighs  $\leq W$ 

Give an algorithm that runs in O(nW) time.



### 0-1 Knapsack

Knapsack 0-1 Knapsack Example Time Complexity

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#### Example

Knapsack 0-1 Knapsack Example Time Complexity

Dynamic Programming

Calculate the best value you can store in a knapsack with W = 7, based on the following price table:

weight w <sub>i</sub>	3	2	1	5	4	
price $p_i$	9	7	3	9	10	

How can we reconstruct the solution (decide which items to include to get the best price)?



#### **Example**

Knapsack Example

Dynamic Programming

Calculate the best value you can store in a knapsack with W = 7, based on the following price table:

 weight  $w_i$  3
 2
 1
 5
 4

 price  $p_i$  9
 7
 3
 9
 10

How can we reconstruct the solution (decide which items to include to get the best price)?



Dynamic Programming

# **Time Complexity**

#### What is the length of the input?

*pseudo-polynomial time:* polynomial if the magnitude of the input numbers is polynomial in the input size

Does this apply to counting sort?



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Summary

## **Dynamic Programming**

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# **Summary of Dynamic Programming**

- optimal substructure: global optimum uses optimal solutions of subproblems
- 2 ordering of subproblems: solve 'smallest' first, build larger solutions from smaller
- 3 'overlapping' subproblems: polynomial number of subproblems, used multiple times
- independent subproblems: optimal solution of one subproblem doesn't affect optimality of another
  - top-down: memoization
  - bottom-up: compute table, then recover solution