Binary Search Trees

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Binary search tree property: a binary tree which is ordered such that every node's left subtree has only smaller values and right subtree has only larger (or equal) values

Useful operations: find, min, max, previous, next, insert, and delete.

Next

If no right child, want minimum ancestor ‘to the right’

NEXT(x)

```
1: if right child exists then
2:     return min under right child
3: else
4:     return UP(x)
```

UP(x)

```
1: p ← x.parent
2: if p = nil or x = p.left then
3:     return p
4: else
5:     return UP(p)
```

Insert

INSERT(n)

```
1: n.parent ← FINDPARENT(n, root, nil)
2: if n.parent is nil then
3:     root ← n
4: else
5:     if n < n.parent then
6:         n.parent.left ← n
7:     else
8:         n.parent.right ← n
```

findParent

FINDPARENT(n, curr, parent)

```
1: if curr is nil then
2:     return parent
3: if n < curr then
4:     return FINDPARENT(n, curr.left, curr)
5: else
6:     return FINDPARENT(n, curr.right, curr)
```

Deletion Outline

3 cases of *delete*(n):

1. no kids: pointer from parent $\leftarrow$ nil
2. 1 child: substitute child for n at parent
3. 2 children: let `next(n)` be s
   a) s takes n's place at parent
   b) n's left subtree becomes s's
   c) somehow rest of n's right subtree becomes s's

let's split 3(c) into 2 cases...

Deletion Outline, Revised

4 cases of $delete(n)$:

1. no kids: pointer from parent $\leftarrow$ nil
2. 1 child: substitute child for n at parent
3. s is n's right child:
   a) substitute s for n
   b) add n's left subtree as s's left subtree
4. s is deeper:
   a) substitute s's right subtree for s
   b) add n's right subtree as s's right subtree
   c) substitute s for n
   d) add n's left subtree as s's left subtree

Moving Subtrees

SUBSTITUTE(old, new)

```
1: if old's parent is nil then
2:     root ← new
3: else
4:     if old is parent's left child then
5:         parent's left child ← new
6:     else
7:         parent's right child ← new
8: if new ≠ nil then
9:     new's parent ← old's parent
```

Deletion

DELETE(n)

```
 1: if n has no left child then
 2:     SUBSTITUTE(n, n's right subtree)
 3: else if n has no right child then
 4:     SUBSTITUTE(n, n's left subtree)
 5: else
 6:     s ← min in n's right subtree
 7:     if n ≠ s.parent then
 8:         SUBSTITUTE(s, s's right subtree)
 9:         s's right subtree ← n's right subtree
10:         s's right child's parent ← s
11:     SUBSTITUTE(n, s)
12:     s's left subtree ← n's left subtree
13:     s's left child's parent ← s
```

Balanced Trees

Structure	Find	Insert	Delete
List (sorted)			
BST (unbalanced)			
BST (balanced)			

Reminder: maintaining a balanced binary tree is important